

## Research Article

# Three-Dimensional SPIHT Coding of Volume Images with Random Access and Resolution Scalability

Emmanuel Christophe<sup>1,2</sup> and William A. Pearlman<sup>3</sup>

<sup>1</sup> Tesa/IRIT 14 port St Etienne, 31500 Toulouse, France

<sup>2</sup> CNES DCT/SI/AP, 18 Avenue E. Belin, 31401 Toulouse, France

<sup>3</sup> Electrical, Computer, and Systems Engineering Department, Rensselaer Polytechnic Institute, Troy, NY 12180-3590, USA

Correspondence should be addressed to Emmanuel Christophe, emmanuel.christophe@cnes.fr

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End users of large volume image datasets are often interested only in certain features that can be identified as quickly as possible. For hyperspectral data, these features could reside only in certain ranges of spectral bands and certain spatial areas of the target. The same holds true for volume medical images for a certain volume region of the subject's anatomy. High spatial resolution may be the ultimate requirement, but in many cases a lower resolution would suffice, especially when rapid acquisition and browsing are essential. This paper presents a major extension of the 3D-SPIHT (set partitioning in hierarchical trees) image compression algorithm that enables random access decoding of any specified region of the image volume at a given spatial resolution and given bit rate from a single codestream. Final spatial and spectral (or axial) resolutions are chosen independently. Because the image wavelet transform is encoded in tree blocks and the bit rates of these tree blocks are minimized through a rate-distortion optimization procedure, the various resolutions and qualities of the images can be extracted while reading a minimum amount of bits from the coded data. The attributes and efficiency of this 3D-SPIHT extension are demonstrated for several medical and hyperspectral images in comparison to the JPEG2000 Multicomponent algorithm.

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## 1. INTRODUCTION

Compression of 3D data volumes poses a challenge to the data compression community. Lossless or near lossless compression is often required for these 3D data, whether medical images or remote sensing hyperspectral images. Due to the huge amount of data involved, even the compressed images are significant in size. In this situation, progressive data encoding enables quick browsing of the image with limited computational or network resources.

For satellite sensors, the trend is toward increase in the spatial resolution, the radiometric precision and possibly the number of spectral bands, leading to a dramatic increase in the amount of bits generated by such sensors. Often, continuous acquisition of data is desired, which requires scan-based mode compression capabilities. Scan-based mode compression denotes the ability to begin the compression of the image when the end of the image is still under acquisition. When the sensor resolution is below one

meter, images containing more than  $30000 \times 30000$  pixels are not exceptional. In these cases, it is important to be able to decode only portions of the whole image. This feature is called random access decoding.

Resolution scalability is another feature that is appreciated within the remote sensing community. Resolution scalability enables the generation of a quick look at the entire image using just few bits of coded data with very limited computation. It also allows the generation of low-resolution images which can be used by applications that do not require fine resolution. More and more applications of remote sensing data are applied within a multiresolution framework [1, 2], often combining data from different sensors. Hyperspectral data should not be an exception to this trend. Hyperspectral data applications are still in their infancy and it is not easy to foresee what the new application requirements will be, but we can expect that these data will be combined with data from other sensors by automated algorithms. Strong transfer constraints are increasingly common in real

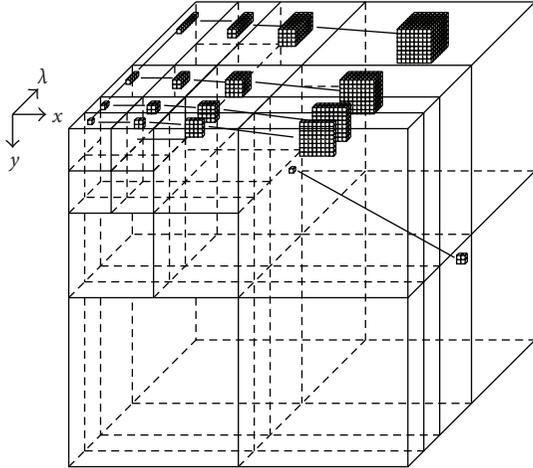


FIGURE 1: Illustration of the wavelet packet decomposition and the tree structure for SPIHT. All descendants for a coefficient  $(i, j, k)$  with  $i$  and  $k$  being odd and  $j$  being even are shown.

remote sensing applications as in the case of the *international charter: space and major disasters* [3]. Resolution scalability is necessary to dramatically reduce the bit rate and provide only the necessary information for the application.

The SPIHT (set partitioning in hierarchical trees) algorithm [4] is a good candidate for onboard hyperspectral data compression. A modified version of SPIHT is currently flying toward the 67P/Churyumov-Gerasimenko comet and is targeted to reach in 2014 (Rosetta mission) among other examples. This modified version of SPIHT is used to compress the hyperspectral data of the VIRTIS instrument [5]. This interest is not restricted to hyperspectral data. The current development of the CCSDS (Consultative Committee for Space Data Systems, which gathers experts from different space agencies as NASA, ESA, and CNES) is oriented toward zerotrees principles [6] because JPEG2000 suffers from implementation difficulties as described in [7] (in the context of implementation compatible with space constraints).

Several papers develop the issue of adaptation from 2D coding to 3D coding using zerotree-based methods. One example is adaptation to multispectral images in [8] through a Karhunen-Loeve transform on the spectral dimension and another is to medical images where [9] uses an adaptation of the 3D SPIHT, first presented in [10]. In [11], a more efficient tree structure is defined and a similar structure proved to be nearly optimal in [12, 13]. To increase the flexibility and the features available as specified in [14], modifications are required. The problem of error resilience is developed in [15] on a block-based version of 3D-SPIHT. A general review of these modifications and a comparison of performances is provided in [16]. Few papers focus on the resolution scalability, as is done in papers [10, 17–20], adapting SPIHT or SPECK (set partitioning-embedded block)[21] algorithms. However, none offers to differentiate the different directions along the coordinate axes to allow full spatial resolution with reduced spectral resolution. In [17, 18], the authors report

a resolution and quality scalable two-dimensional SPIHT, but without the random access capability to be enabled in our proposed algorithm. Our proposed extension to three dimensions with random access decodability that retains spatial and quality scalability requires significant changes of the transform and tree structure and search mode, and the addition of a post-compression rate allocation procedure. To the authors' knowledge, no previous work presents the combination of all these features doing a rate distortion optimization between blocks, while maintaining optimal rate-distortion performance and preserving the properties of spatial and quality scalability.

This paper presents the extension of the well-known SPIHT algorithm for 3D data enabling random access and resolution scalability, while keeping quality and rate scalability and extends the previous work presented in [22]. Compression performance and attributes are compared with JPEG2000 [23].

## 2. DATA DECORRELATION AND TREE STRUCTURE

### 2.1. 3D anisotropic wavelet transform

Hyperspectral images contain one image of the scene for different wavelengths, thus two dimensions of the 3D hyperspectral cube are spatial and the third one is spectral (in the wavelength ( $\lambda$ ) sense). Medical magnetic resonance (MR) or computed tomography (CT) images contain one image for each slice of observation, in which case the three dimensions are spatial. However, the resolution and statistical properties of the third direction are different. To avoid confusion, the first two dimensions are referred to as spatial, whereas the third one is called spectral. An anisotropic 3D wavelet transform is applied to the data for the decorrelation. This decomposition consists of performing a classic dyadic 2D wavelet decomposition on each image followed by a 1D dyadic wavelet decomposition in the third direction. The obtained subband organization is represented in Figure 1 and is also known as wavelet packet. The decomposition is nonisotropic as not all subbands are regular cubes and some directions are privileged. It has been shown that this anisotropic decomposition is nearly optimal in a rate-distortion sense in terms of entropy [24] as well as real coding [12]. To the authors' knowledge, this is valid for 3D hyperspectral data as well as for 3D magnetic resonance medical images and video sequences. Moreover, this is the only 3D wavelet transform supported by the JPEG2000 standard in Part II [25].

### 2.2. Tree structure

The SPIHT algorithm [4] uses a tree structure to define a relationship between wavelet coefficients from different subbands. To adapt the SPIHT algorithm on the anisotropic 3-D (wavelet packet) decomposition, a suitable tree structure must be defined. Let us define  $\mathcal{O}_{\text{spat}}(i, j, k)$  as the spatial ( $x - y$  band in Figure 1) offspring of the pixel located at sample  $i$ , line  $j$  in band  $k$ . The first coefficient in the upper front, left corner is noted as  $(0, 0, 0)$ . In the spatial direction,

the relation is similar to the one defined in the original SPIHT. In general, we have  $\mathcal{O}_{\text{spat}}(i, j, k) = \{(2i, 2j, k), (2i + 1, 2j, k), (2i, 2j + 1, k), (2i + 1, 2j + 1, k)\}$ . In the highest spatial frequency subbands, there are no offspring:  $\mathcal{O}_{\text{spat}}(i, j, k) = \emptyset$ . In the lowest frequency subband, coefficients are grouped in  $2 \times 2$  as in the original SPIHT. Let  $n_s$  denote the number of samples per line and  $n_l$  the number of lines in the lowest frequency subband. We have for  $(i, j, k)$  in the lowest frequency subband:

- (i) if  $i$  even and  $j$  even:  $\mathcal{O}_{\text{spat}}(i, j, k) = \emptyset$ ;
- (ii) if  $i$  odd and  $j$  even:  $\mathcal{O}_{\text{spat}}(i, j, k) = \{(i + n_s - 1, j, k), (i + n_s, j, k), (i + n_s - 1, j + 1, k), (i + n_s, j + 1, k)\}$ ;
- (iii) if  $i$  even and  $j$  odd:  $\mathcal{O}_{\text{spat}}(i, j, k) = \{(i, j + n_l - 1, k), (i + 1, j + n_l - 1, k), (i, j + n_l, k), (i + 1, j + n_l, k)\}$ ;
- (iv) if  $i$  odd and  $j$  odd:  $\mathcal{O}_{\text{spat}}(i, j, k) = \{(i + n_s - 1, j + n_l - 1, k), (i + n_s, j + n_l - 1, k), (i + n_s - 1, j + n_l, k), (i + n_s, j + n_l, k)\}$ .

The spectral ( $\lambda$  direction in Figure 1) offspring  $\mathcal{O}_{\text{spec}}(i, j, k)$  are defined in a similar way, but only for the lowest spatial subband: if  $i \geq n_s$  or  $j \geq n_l$ , we have  $\mathcal{O}_{\text{spec}}(i, j, k) = \emptyset$ . Otherwise, apart from the highest and lowest spectral frequency subbands, we have  $\mathcal{O}_{\text{spec}}(i, j, k) = \{(i, j, 2k), (i, j, 2k + 1)\}$  for  $i < n_s$  and  $j < n_l$ . In the highest spectral frequency subbands, there are no offspring:  $\mathcal{O}_{\text{spec}}(i, j, k) = \emptyset$  and in the lowest, coefficients are grouped by 2 to have a construction similar to SPIHT. Let  $n_b$  be the number of spectral bands in the lowest spectral frequency subband:

- (i) if  $i < n_s, j < n_l, k$  even:  $\mathcal{O}_{\text{spec}}(i, j, k) = \emptyset$ ;
- (ii) if  $i < n_s, j < n_l, k$  odd:  $\mathcal{O}_{\text{spec}}(i, j, k) = \{(i, j, k + n_b - 1), (i, j, k + n_b)\}$ .

With these relations, we have a separation in nonoverlapping trees of all the coefficients of the wavelet transform of the image. The tree structure is illustrated in Figure 1 for three levels of decomposition in each direction. Each of the coefficients is the descendant of a root coefficient located in the lowest frequency subband. It has to be noted that all the coefficients belonging to the same tree correspond to a geometrically similar area of the original image, in the three dimensions.

We can compute the maximum number of coefficients in a tree rooted at  $(i, j, k)$  for a 5 level spatial and spectral decomposition. The maximum of descendants occurs when  $k$  is odd and at least either  $i$  or  $j$  is odd. For this situation, we have  $1 + 2 + 2^2 + \dots + 2^5 = 2^6 - 1$  spectral descendants (including the root) and for each of these, we have  $1 + 2^2 + (2^2)^2 + (2^3)^2 + \dots + (2^5)^2 = 2^0 + 2^2 + 2^4 + \dots + 2^{10} = (2^{12} - 1)/3$  spatially linked coefficients. Let  $l_{\text{spec}}$  be the number of decompositions in the spectral direction and let  $l_{\text{spac}}$  be the same in the spatial direction, we obtain the general formula:

$$n_{\text{desc}} = (2^{l_{\text{spac}}+1} - 1) \frac{2^{2(l_{\text{spec}}+1)} - 1}{3} \quad (1)$$

for the maximum number of coefficients in a tree. Thus the number of coefficients in the tree is at most 85995 ( $l_{\text{spec}} = 5$  and  $l_{\text{spac}} = 5$ ) if the given coefficient has both spectral and

spatial descendants. Coefficient  $(0, 0, 0)$ , for example, has no descendants at all.

### 3. BLOCK CODING

To provide random access, it is necessary to encode separately different areas of the image. Encoding separately portions of the image provides several advantages. First, scan-based mode compression is made possible as the whole image is not necessary. Once again, we do not consider here the problem of the scan-based wavelet transform which is a separate issue. Secondly, encoding parts of the image separately also provides the ability to use different compression parameters for different parts of the image, enabling the possibility of high-quality region of interest (ROI) and the possibility of discarding unused portions of the image. An unused portion of the image could be an area with clouds in remote sensing or irrelevant organs in a medical image. Third, transmission errors have a more limited effect in the context of separate coding; the error only affects a limited portion of the image.

Direct transformation and coding different portions of the image results in poor coding efficiency and blocking artifacts visible at boundaries between adjacent portions. However, if we encode portions of the full-image transform corresponding to image regions that together constitute the whole, coding efficiency is maintained and boundary artifacts vanish in the inverse transform process. This strategy has been used for this particular purpose on the EZW algorithm in [26], and in [15] for 3D-SPIHT in the context of video coding. Finally, one limiting factor of the SPIHT algorithm is the complicated list processing requiring a large amount of memory. If the processing is done only on one part of the transform at a time, the number of coefficients involved is dramatically reduced and so is the memory necessary to store the control lists in SPIHT.

With the tree structure defined in Section 2, a natural block organization appears. A tree-block (later simply referred to as *block*) is generated by 8 coefficients forming a  $2 \times 2 \times 2$  cube from the lowest subband together with all their descendants. It is more easily visualized in the two-dimensional case in Figure 2, where is shown a  $2 \times 2$  group in the lowest frequency subband and all its descendants forming a tree-block. All the coefficients linked to the root coefficient in the lowest subband shown for three dimensions on Figure 1 are part of the same tree-block together with seven other trees. Grouping the coefficients by 8 enables the use of neighbor similarities between coefficients. This grouping of coefficients in the lowest frequency subband is analogous to the grouping of  $2 \times 2$  in the original SPIHT patent [27] and paper [4]. The gray-shaded coefficients in Figure 1 constitute a block in our three-dimensional transform.

For this grouping, the number of coefficients in each block will be the same, the only exception being the case where at least one dimension of the lowest subband is odd. In a  $2 \times 2 \times 2$  root group, we have three coefficients which have the full sets of descendants, whose number is given by (1), three have only spatial descendants, one has only spectral descendants, and the last one has no descendant. The number of coefficients in a block, which

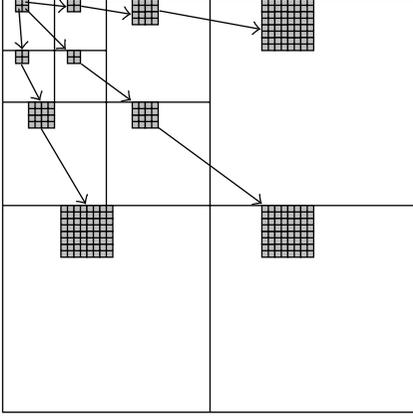


FIGURE 2: Equivalence of the block structure for 2D, all coefficients in gray belong to the same block. In the following algorithm, an equivalent 3D block structure is used.

determines the maximum amount of memory necessary for the compression, will finally be  $262144 = 2^{18}$  (valid for 5 decompositions in the spatial and spectral directions).

The granularity of the random access obtained with this method is very small. Spatially, the grain size is  $2 \times 2$ , compared to JPEG2000's  $32 \times 32$  or  $64 \times 64$ , which are the typical sizes of the encoded subblocks of the subbands. Using subblocks smaller than  $32 \times 32$  in JPEG2000 results in considerable loss of coding efficiency. JPEG2000 encodes the spectrally transformed slices or spectral bands independently, so its grain size in the spectral direction is 1 versus 2 for our method. With the  $2 \times 2 \times 2$  root group, it is possible to retrieve almost only the required coefficients to decode a given area. Moreover, every coefficient can be retrieved only to the bit plane necessary to give the expected quality.

## 4. ENABLING RESOLUTION SCALABILITY

### 4.1. Original SPIHT algorithm

The original SPIHT algorithm processes the coefficients bit plane by bit plane. Coefficients are stored in three different lists according to their significance. The list of significant pixels (LSP) stores the coefficients that have been found significant in a previous bit plane and that will be refined in the following bit planes. Once a coefficient is on the LSP, it remains significant at all lower thresholds. It stays on the LSP, so that it can be successively refined with bits from its lower bit planes. The list of insignificant pixels (LIP) contains the coefficients which are still insignificant, relative to the current bit plane and which are not part of a tree from the third list (LIS). Coefficients in the LIP are transferred to the LSP when they become significant. The third list is the list of insignificant sets (LIS). A set is said to be insignificant if all descendants, in the sense of the previously defined tree structure, are not significant in the current bit plane. For the

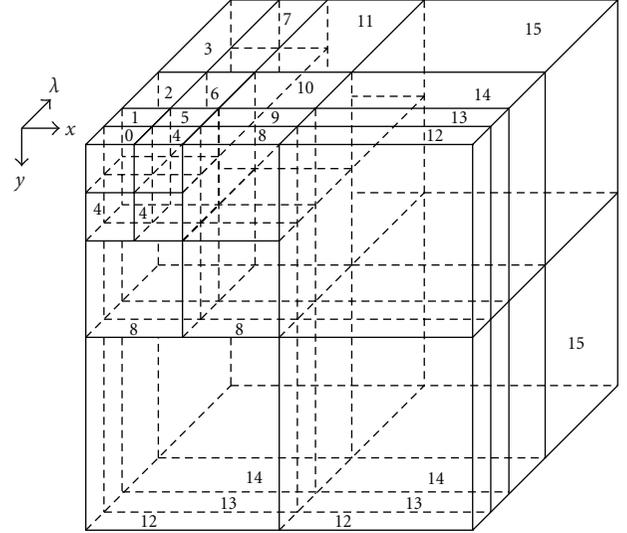


FIGURE 3: Illustration of the resolution level numbering. If a low-resolution image is required (either spectral or spatial), only subbands with a resolution number corresponding to the requirements are processed.

bit plane  $t$ , we define the significance function  $S_t$  of a set  $\mathcal{T}$  of coefficients :

$$S_t(\mathcal{T}) = \begin{cases} 0 & \text{if } \forall c \in \mathcal{T}, |c| < 2^t \\ 1 & \text{if } \exists c \in \mathcal{T}, |c| \geq 2^t. \end{cases} \quad (2)$$

If  $\mathcal{T}$  consists of a single coefficient, we denote its significance function by  $S_t(i, j, k)$ .

Let  $\mathcal{D}(i, j, k)$  be all descendants of  $(i, j, k)$ ,  $\mathcal{O}(i, j, k)$  only the offspring (i.e., the first-level descendants) and  $\mathcal{L}(i, j, k) = \mathcal{D}(i, j, k) - \mathcal{O}(i, j, k)$ , the granddescendant set. A type A tree is a tree where  $\mathcal{D}(i, j, k)$  is insignificant (all descendants of  $(i, j, k)$  are insignificant); a type B tree is a tree where  $\mathcal{L}(i, j, k)$  is insignificant (all granddescendants of  $(i, j, k)$  are insignificant). The full SPIHT algorithm can be found in [4].

### 4.2. Introducing resolution scalability

In SPIHT, there is no distinction between coefficients from different resolution levels. To provide resolution scalability, we need to provide the ability to decode only the coefficients from a selected resolution. A resolution comprises 1 or 3 subbands. To enable this capability, we keep three lists for each resolution level  $r$ . When  $r = 0$ , only coefficients from the low-frequency subbands will be processed. Resolution levels must be processed in increasing order because to reconstruct a given resolution, all the lower-order resolution levels are needed. Coefficients are processed according to the resolution level to which they correspond. For a 5-level wavelet decomposition in the spectral and spatial direction, a total of 36 resolution levels will be available (illustrated on Figure 3 for 3-level wavelet and 16 resolution levels available). Each level  $r$  keeps in memory three lists: LSP<sub>r</sub>, LIP<sub>r</sub>, and LIS<sub>r</sub>.

Some difficulties arise from this organization and the progression order to follow; several options are available (Figure 4). If the priority is given to full-resolution scalability compared to the bit plane scalability, some extra precautions have to be taken. The different possibilities for scalability order are discussed in Section 4.3. In the most complicated case, where all bit planes for a given resolution  $r$  are processed before the descendant resolution  $r_d$  (full-resolution scalability), the last element to process for  $LSP_{r_d}$ ,  $LIP_{r_d}$ , and  $LIS_{r_d}$  for each bit plane  $t$  has to be remembered. Details of the resolution-scalable algorithm, referred as SPIHT RARS (Random Access with Resolution Scalability) are given in Algorithm 1.

This new algorithm, which processes all bit planes at a given resolution level, provides strictly the same code bits as the original SPIHT. The bits are just organized in a different order. With the block structure, memory footprint during compression is dramatically reduced. The resolution scalability with its several lists does not increase the amount of memory necessary as the coefficients are just spread onto different lists.

### 4.3. Switching loops

The priority of scalability type can be chosen by the progression order of the two “for” loops (just after the initialization stage) in the 3D SPIHT RARS algorithm. As written, the priority is resolution scalability, but these loops can be inverted to give priority to quality scalability. The different progression orders are illustrated in Figures 4(a) and 4(b). Processing the resolution completely before proceeding to the next one (Figure 4(b)) requires more precautions.

When processing resolution  $r$ , a significant descendant set is partitioned into its offspring in  $r_d$  and its granddescendant set. Therefore, some coefficients are added to  $LSP_{r_d}$  in the step marked (2) in the algorithm (similar for the  $LIP_{r_d}$  and  $LIS_{r_d}$ ). This is an additional step compared to the original SPIHT [4]. So even before processing resolution  $r_d$ , the  $LSP_{r_d}$  may contain some coefficients which were added at different bit planes. One possible content of an  $LSP_{r_d}$  could be

$$LSP_{r_d} = \{(i_0, j_0, k_0)(t_{19}), (i_1, j_1, k_1)(t_{19}), \dots, \\ (i_n, j_n, k_n)(t_{12}), \dots, \\ (i_{n'}, j_{n'}, k_{n'})(t_0), \dots\}, \quad (3)$$

(the bit plane when a coefficient was added to the list is given in parentheses following the coordinate) 19 being the highest bit plane in this case (depending on the image).

When we process  $LSP_{r_d}$ , we should skip entries added at lower bit planes than the current one. For example, there is no meaning to refine a coefficient added at  $t_{12}$  when we are working in bit plane  $t_{18}$ .

Furthermore, at the step marked (1) in the algorithm above, when processing resolution  $r_d$  we add some coefficients to  $LSP_{r_d}$ . These coefficients have to be added at the proper position within  $LSP_{r_d}$  to preserve the order. When adding a coefficient at step (1) for the bit plane  $t_{19}$ , we insert it just after the other coefficient from bit plane  $t_{19}$  (at the end of

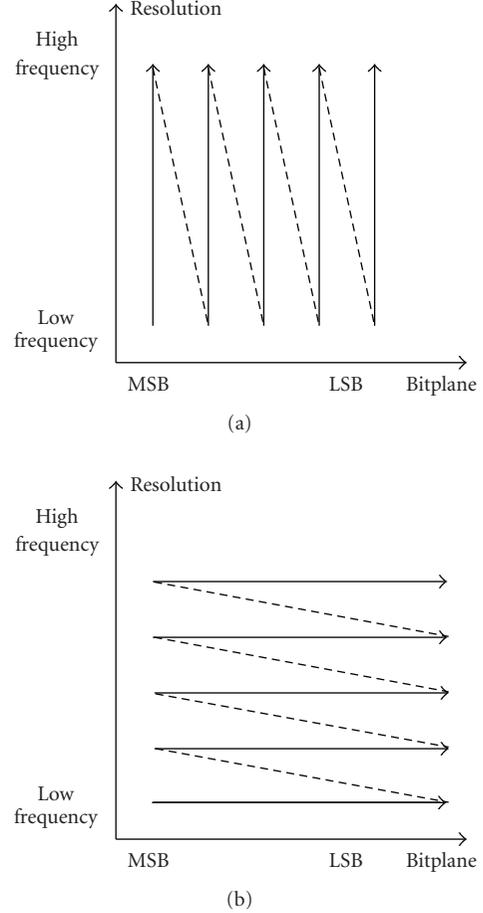


FIGURE 4: Scanning order for SNR scalability (a) or resolution scalability (b).

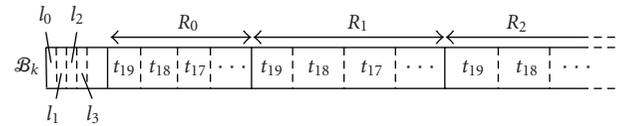


FIGURE 5: Resolution scalable bitstream structure with header. The header allows the decoder to jump directly to resolution 1 without completely decoding or reading resolution 0.  $R_0, R_1, \dots$  denote the different resolutions,  $t_{19}, t_{18}, \dots$  the different bit planes.  $l_i$  is the size in bits of  $R_i$ .

the first line of (3). Keeping the order avoids looking through the whole list to find the coefficients to process at a given bit plane and can be done simply with a pointer.

The bitstream structure obtained for this algorithm is shown in Figure 5 and is called the resolution scalable structure. If quality scalability replaces resolution scalability as a priority, the “for” loops, that step through resolutions and bit planes, can be inverted to process one bit plane completely for all resolutions before going to the next bit plane. In this case, the bitstream structure obtained is different and illustrated in Figure 6 and is called the quality

```

// Initialization step:
t ← number of bit planes
LSP0 ← ∅
LIP0 ← all the coefficients without any parents (the 8 root
coefficients of the block);
LIS0 ← all coefficients from the LIP0 with descendants (7
coefficients as only one has no descendant);
For r ≠ 0, LSPr ← ∅, LIPr ← ∅, LISr ← ∅;
// List processing:
for each r from 0 to maximum resolution do
  for each t from the highest bit plane to 0 (bit planes) do
    // Sorting pass:
    for each entry (i, j, k) of the LIPr which had been
      added at a threshold strictly greater to the current t do
      Output St(i, j, k);
      If St(i, j, k) = 1, move (i, j, k) to LSPr and
      output the sign of ci,j,k (1);
    for each entry (i, j, k) of the LISr which had been
      added at a threshold greater or equal to the current t
      do
      if the entry is type A then
        Output St(D(i, j, k));
        if St(D(i, j, k)) = 1 then
          for all (i', j', k') ∈ O(i, j, k) do
            output St(i', j', k');
            if St(i', j', k') = 1 then
              add (i', j', k') to the LSPrd;
              Output the sign of ci',j',k';
            else add (i', j', k') to the end of the
              LIPrd (2);
          if L(i, j, k) ≠ ∅ then move (i, j, k) to the
            LISr as a type B entry;
          else remove (i, j, k) from the LISr;
      if the entry is type B then
        Output St(L(i, j, k));
        if St(L(i, j, k)) = 1 then
          Add all the (i', j', k') ∈ O(i, j, k) to the
            LISrd as a type A entry;
          Remove (i, j, k) from the LISr;
    //Refinement pass:
    for all entries (i, j, k) of the LSPr which had been
      added at a threshold strictly greater than the current
      t do
      Output the tth most significant bit of ci,j,k

```

ALGORITHM 1: Resolution scalable 3D SPIHT RARS.

scalable structure. The differences between scanning order are shown in Figure 4.

#### 4.4. More flexibility at decoding

3D SPIHT RARS possesses great flexibility and the same image can be encoded up to an arbitrary resolution level or down to a certain bit plane, depending on the two possible loop orders. The decoder can just proceed to the same level to decode the image. However, an interesting feature to have is the possibility to encode the image only once, with all resolution and all bit planes and then during the decoding to

choose which resolution and which bit plane to decode. One may need only a low-resolution image with high-radiometric precision or a high-resolution portion of the image with rough-radiometric precision.

When the resolution scalable structure is used (Figure 5), it is easy to decode up to the desired resolution, but if not all bit planes are necessary, we need a way to jump to the beginning of resolution 1 once resolution 0 is decoded for the necessary bit planes. The problem is the same with the quality scalable structure (Figure 6) exchanging bit plane and resolution in the problem description.

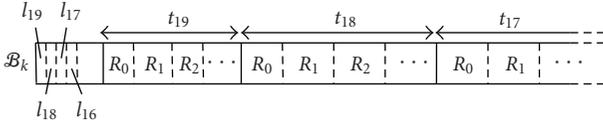


FIGURE 6: Quality scalable bitstream structure with header. The header allows the decoder to continue the decoding of a lower bit plane without having to finish all the resolution at the current bit plane.  $R_0, R_1, \dots$  denote the different resolutions,  $t_{19}, t_{18}, \dots$  the different bit planes.  $l_i$  is the size in bits of the bit plane corresponding to  $t_i$ .

To overcome this problem, we need to introduce a block header describing the size of each portion of the bitstream. The cost of this header is negligible: the number of bits for each portion is coded with 24 bits, enough to code part sizes up to 16 Mbits. The lowest resolutions (resp., the highest bit planes) which are using only few bits will be processed fully, regardless of the specification at the decoder, as the cost in size and processing is low and therefore their sizes need not to be kept. Only the sizes of long parts are kept: we do not keep the size individually for the first few bit planes or the first few resolutions, since they will be decoded in any case. Only the sizes of lower bit planes and higher resolutions (in general well above 10000 bits), which comprise about 10 numbers (each coded with 32 bits to allow sizes up to 4 Gb), need to be written to the bitstream. Then this header cost will remain below 0.1%.

As in [17], simple markers could have been used to identify the beginning of new resolutions of new bit planes. Markers have the advantage to be shorter than a header coding the full size of the following block. However, markers make the full reading of the bitstream compulsory and the decoder cannot just jump to the desired part. As the cost of coding the header remains low, this solution is chosen.

## 5. DRAWBACKS OF BLOCK PROCESSING AND INTRODUCTION OF RATE ALLOCATION

### 5.1. Rate allocation and keeping the SNR scalability

The problem of processing different areas of the image separately always resides in the rate allocation for each of these areas. A fixed rate for each area is usually not a suitable decision as complexity most probably varies across the image. If quality scalability is necessary for the full image, we need to provide the most significant bits for one block before finishing the previous one. This could be obtained by cutting the bitstream for all blocks and interleaving the parts in the proper order. With this solution, the rate allocation will not be available at the bit level due to the block organization and the spatial separation, but a tradeoff with quality layers organization can be used.

### 5.2. Layer organization and rate-distortion optimization

The idea of quality layers is to provide different targeted bit rates in the same bitstream [28]. For example, a bitstream can

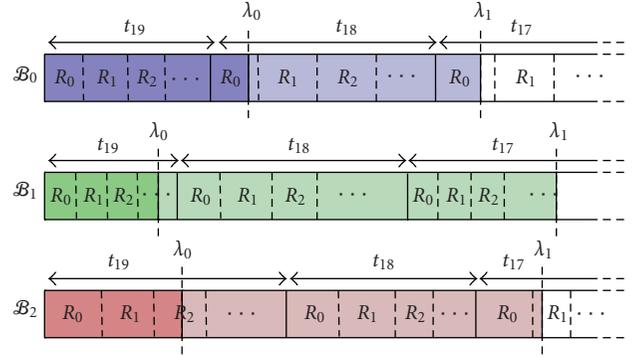


FIGURE 7: An embedded scalable bitstream generated for each block  $\mathcal{B}_k$ . The rate-distortion algorithm selects different cutting points corresponding to different values of the parameter  $\lambda$ . The final bitstream is illustrated in Figure 8.

provide two quality layers: one at 1.0 bits per pixel (bpp) and another at 2.0 bpp. If the decoder needs a 1.0 bpp image, just the beginning of the bitstream is transferred and decoded. If a higher-quality image is needed, the first layer is transmitted, decoded, and then refined with the information from the second layer.

As the bitstream for each block is already embedded, to construct these layers, we just need to select the cutting points for each block and each layer leading to the correct bit rate with the optimal quality for the entire image. Once again, it has to be a global optimization and not only local, as complexity will vary across blocks.

A simple Lagrangian optimization method [29] gives the optimal cutting point for each block  $\mathcal{B}_k$ . This optimization consists in minimizing the cost function  $J(\lambda) = \sum_k (D_k + \lambda R_k)$ :  $D_k$  being the distortion of the block  $\mathcal{B}_k$ ,  $R_k$  its rate, and  $\lambda$  the Lagrange parameter. This Lagrangian optimization to find the cutting point between different blocks is also used in JPEG2000 and referred to as PCRDOpt (post-compression rate-distortion optimization) [28]. It has to be noted that the progressive bit plane coding of SPIHT provides a straightforward implementation of this method.

The result of the Lagrangian optimization led to an interleaved bitstream between different blocks, as described in Figures 7 and 8.

### 5.3. Low-cost distortion tracking: during the compression

In the previous part, we assumed that the distortion was known for every cutting point (every bit in fact) of the bitstream for one block. As the bitstream for one block is in general about millions of bits, it is too costly to keep all this distortion information in memory. Only a few hundred cutting points are recorded with their rate and distortion information.

Getting the rate for one cutting point is the easy part: one just has to count the number of bits before this point. The distortion requires more processing. The distortion value during the encoding of one block can be obtained

TABLE 1: Data sets.

Image	Type	Dynamic	Size	$\sigma^2$
Moffett sc1	Hyperspectral	16 bits	$512 \times 512 \times 224$	1749626.1
Moffett sc3	Hyperspectral	16 bits	$512 \times 512 \times 224$	1666647.5
Jasper sc1	Hyperspectral	16 bits	$512 \times 512 \times 224$	1361347.8
Cuprite sc1	Hyperspectral	16 bits	$512 \times 512 \times 224$	3212383.2
CT_skull	CT	8 bits	$256 \times 256 \times 192$	3201.0540
CT_wrist	CT	8 bits	$256 \times 256 \times 176$	2431.6957
MR_sag_head	MR	8 bits	$256 \times 256 \times 56$	671.58119
MR_ped_chest	MR	8 bits	$256 \times 256 \times 64$	444.34858

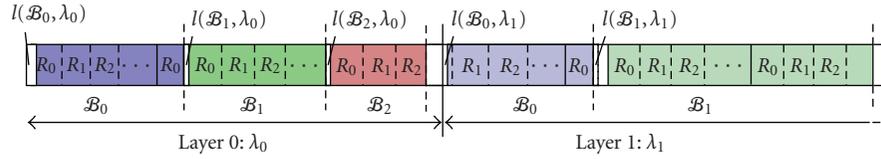


FIGURE 8: The bitstreams are interleaved for different quality layers. To permit the random access to the different blocks, the length in bits of each part corresponding to a block  $\mathcal{B}_k$  and a quality layer corresponding to  $\lambda_q$  is given by  $l(\mathcal{B}_k, \lambda_q)$ .

with a simple tracking. Let us consider the instant in the compression when the encoder is adding one precision bit for one coefficient  $c$  at the bit plane  $t$ . Let  $c_t$  denote the new approximation of  $c$  in the bit plane  $t$  given by adding this new bit.  $c_{t+1}$  was the approximation of  $c$  at the previous bit plane.

SPIHT uses a deadzone quantizer, so if the refinement bit is 0, we have  $c_t = c_{t+1} - 2^{t-1}$  and if the refinement bit is 1, we have  $c_t = c_{t+1} + 2^{t-1}$ . Let us call  $D^a$  the total distortion of the block after this bit was added and  $D^b$  the total distortion before. We have the following:

- (i) with a refinement bit of 0:

$$\begin{aligned}
 D^a - D^b &= (c - c_t)^2 - (c - c_{t+1})^2 \\
 &= (c_{t+1} - c_t)(2c - c_t - c_{t+1}) \\
 &= 2^{t-1}(2(c - c_{t+1}) + 2^{t-1}),
 \end{aligned} \tag{4}$$

giving

$$D^a = D^b + 2^{t-1}(2(c - c_{t+1}) + 2^{t-1}); \tag{5}$$

- (ii) with a refinement bit of 1:

$$D^a = D^b - 2^{t-1}(2(c - c_{t+1}) - 2^{t-1}). \tag{6}$$

Since this computation can be done using only right and left bit shifts and additions, the computational cost is low. The algorithm does not need to know the initial distortion value as the rate-distortion method holds if *distortion* is replaced by *distortion reduction*. The value can be high and has to be kept internally in a 64-bit integer. As seen before, we have  $2^{18}$  coefficients in one block, and for some of them, the value can reach  $2^{20}$ . Therefore, 64 bits seem to be a reasonable choice that remains valid for the worst cases.

The evaluation of the distortion is done in the transform domain, directly on the wavelet coefficients. This can be

done only if the transform is orthogonal. The 9/7 transform is approximately orthogonal. In [30], the computation of the weight to apply to each wavelet subband for the rate allocation is detailed. The weight can be introduced as in (5) and (6) as a multiplicative factor to get a precise distortion evaluation in the wavelet domain. However, the gain in quality introduced by the increase in precision is negligible (about 0.01 dB) compared to the increase in complexity. Thus these weights are not kept in the following results.

## 6. RESULTS

### 6.1. Data and performance measurement

The hyperspectral data subsets originate from the airborne visible infrared imaging spectrometer (AVIRIS) sensor. We use radiance unprocessed data. The original AVIRIS scenes are  $614 \times 512 \times 224$  pixels. For the simulations here, we crop the data to  $512 \times 512 \times 224$  starting from the upper left corner of the scene. To make comparison easier with other papers, we use well-known data sets: particularly scenes 1 and 3 of the run from AVIRIS on Moffett Field, but also scene 1 over Jasper Ridge and scene 1 over Cuprite site. MR and CT medical images are also used. The details of all the images are given in Table 1.

Error is given in terms of signal-to-noise ratio (SNR), root mean square error (RMSE), and maximum error  $e_{\max}$ . SNR is computed according to the variance ( $\sigma^2$ ) values from Table 1:  $\text{SNR} = 10\log_{10}\sigma^2/\text{MSE}$ . All errors are measured in the final reconstructed dataset compared to the original data. Choosing a distortion measure suitable to hyperspectral data is not easy matter as shown in [31]. The rate-distortion optimization is based on the additive property of the distortion measure and optimized for the mean squared error (MSE). Our goal here is to choose an acceptable

distortion measure for general use on different kinds of volume images. The MSE-based distortion measures here are appropriate and popular and are selected to facilitate comparisons.

Final rate is calculated directly from the size of the codestream and includes all headers and required side information. This rate is given in terms of bits per pixel per band (bpppb), where *band* means spectral band for hyperspectral data and axial slice for medical data.

An optional context-based arithmetic coder is included to improve rate performance [32]. In the context of a reduced complexity algorithm, the slight improvement in performance introduced by the arithmetic coder does not seem worth the complexity increase. Results with arithmetic coder are given for reference in Table 2. Unless stated otherwise, results in this paper do not include the arithmetic coder. Several particularities have to be taken into account to preserve the bitstream flexibility. First, contexts of the arithmetic coder have to be reset at the beginning of each part to be able to decode the bitstream partially. Secondly, the rate recorded during the rate-distortion optimization has to be the rate provided by the arithmetic coder.

The raw compression performances of the previously defined random access with resolution scalability (3D-SPIHT-RARS) are compared with the best up to date method without taking into account the specific properties available for the previously defined algorithm. The reference results are obtained with the version 5.0 of Kakadu software [33] using the JPEG2000 Part 2 options: wavelet intercomponent transform to obtain a transform similar to the one used by our algorithm. SNR values are similar to the best values published in [34]. The results were also confirmed using the latest reference implementation of JPEG2000, the *verification model* (VM) version 9.1. Our results are not expected to be better, but are here to show that the increase in flexibility does not come with a prohibitive cost in performance. It also has to be noted that the results presented here for 3D-SPIHT-RARS do not include any entropy coding of the SPIHT sorting output, thus simplifying the implementation.

## 6.2. Performance comparisons

First, coding results are compared with the original SPIHT. The decrease in quality is very low at 1 bpppb (under 0.05 dB) and remain low at 0.5 bpppb (about 0.40 dB). The source of performance decrease is the separation of the wavelet subbands at each bit plane which causes different bits to be kept if the bitstream is truncated. Once again, if lossless compression is required, the two algorithms, SPIHT and SPIHT-RARS, provide exactly the same bits reordered (apart from the headers).

Computational complexity is not easy to measure, but one way to get a rough estimation is to measure the time needed for the compression of one image. The version of 3D-SPIHT here is a demonstration version and there is a lot of room for improvement. The compression time with similar options is 20 s for Kakadu v5.0, 600 s for VM 9.1, and 130 s for 3D-SPIHT-RARS. These values are given only to show that compression time is reasonable for a demonstration

TABLE 2: Lossless compression rates (bpppb) (results denoted with \* use the additional lifting steps from [9]).

Image	JPEG2000 MT	SPIHT-RARS	SPIHT-RARS (with AC)
CT_skull	2.93	2.21	<b>2.16</b>
CT_wrist	1.78	1.30	<b>1.27</b>
MR_sag_head	<b>2.30</b>	2.41	2.35
MR_ped_chest	2.00	1.96	<b>1.92</b>
Moffett sc3	<b>5.14</b>	5.37*	5.29*
Moffett sc1	<b>5.65</b>	5.83*	5.75*
Jasper sc1	<b>5.55</b>	5.74*	5.67*
Cuprite sc1	<b>5.29</b>	5.51*	5.43*

TABLE 3: Quality for different rates for Moffett sc3.

	Rate (bpppb)	2.0	1.0	0.5	0.1
SNR	JPEG2000 MT	54.90	48.63	43.52	33.16
	3D-SPIHT-RARS	54.07	47.84	42.49	32.28
RMSE	JPEG2000 MT	2.32	4.78	8.61	28.39
	3D-SPIHT-RARS	2.56	5.24	9.69	31.42
$e_{\max}$	JPEG2000 MT	24	66	157	1085
	3D-SPIHT-RARS	37	80	161	1020

implementation and the comparison with the demonstration implementation of JPEG2000, VM9.1 shows that this is the case. The value given here for 3D-SPIHT-RARS includes the 30 seconds necessary to perform the 3D wavelet transform with QccPack.

Table 2 compares the lossless performance of the two algorithms. JPEG2000 is used with a multicomponent transform (MT). For both, the same integer 5/3 wavelet transform is performed with the same number of decompositions in each direction. The modified 5/3 wavelet with additional lifting steps from [9] is also compared.

Performances between the algorithms are quite similar for the MR images. SPIHT-RARS outperforms JPEG2000 on the CT images, but JPEG2000 gives a lower bit rate for hyperspectral images. It has to be noted that the original 5/3 wavelet transform gives better results for the medical images while the modified transform performs better on hyperspectral images.

Table 3 compares the lossy performances of the two algorithms in terms of different quality criteria and Table 4 provides the SNR obtained on several popular datasets to facilitate comparisons. It is confirmed that the increase in flexibility of the 3D-SPIHT-RARS algorithm does not come with a prohibitive impact on performances. We can observe less than 1 dB difference between the two algorithms. A noncontextual arithmetic coder applied directly on the 3D-SPIHT-RARS bitstream already reduces this difference to 0.4 dB (not used in the presented results).

## 6.3. Resolution scalability from a single bitstream

Different resolutions and different quality levels can be retrieved from one bitstream. Table 5 presents different

TABLE 4: SNR for popular data sets.

	Rate (bpppb)	2.0	1.0	0.5	0.1
Moffett sc1	JPEG2000 MT	51.99	45.48	40.18	29.75
	3D-SPIHT-RARS	50.87	44.27	39.32	28.82
Jasper sc1	JPEG2000 MT	51.60	44.85	39.31	28.99
	3D-SPIHT-RARS	50.52	43.55	38.34	28.03
Cuprite sc1	JPEG2000 MT	56.72	50.99	46.71	38.72
	3D-SPIHT-RARS	55.35	50.15	46.04	37.98
CT_skull <sup>(a)</sup>	JPEG2000 MT	—	39.15	33.23	23.02
	3D-SPIHT-RARS	—	37.63	31.92	22.39
MR_sag_head <sup>(b)</sup>	JPEG2000 MT	—	33.08	27.53	19.04
	3D-SPIHT-RARS	—	31.91	26.60	18.45

<sup>(a)</sup> PSNR values can be obtained by adding 13.08

<sup>(b)</sup> PSNR values can be obtained by adding 19.85

results on Moffett Field scene 3 changing the number of resolutions and bit planes to decode the bitstream. The compression is done only once and the final bitstream is organized in different parts corresponding to different resolution and quality. From this single-compressed bitstream, all these results are obtained by changing the decoding parameters. Different bit depths and different resolutions are chosen arbitrarily to obtain a lower resolution and lower quality image. Distortion measures are provided for the lower resolution image as well as the bit rate necessary to transmit or store this image.

For the results presented in Table 5, similar resolutions are chosen for spectral and spatial directions, but this is not mandatory as illustrated in Figure 9. The reference low-resolution image is the low-frequency subband of the wavelet transform up to the desired level. To provide an accurate radiance value, coefficients are scaled properly to compensate gains due to the wavelet filters (depending on the resolution level).

Table 5 shows, for example, that discarding the 6 lower bit planes, a half resolution image can be obtained with a bit rate of 0.203 bpppb and an RMSE of 6.47 (for this resolution).

We can see that at high quality, decoding to lower resolution greatly decreases the retrieval time. An algorithm working with hyperspectral data could choose to discard 4 bit planes and to work at 1/4 resolution, thereby reducing the amount of data to process by a factor of 10, and enabling simple onboard processing while keeping a good spectral quality (detection of area of interest, detect the clouds to discard useless information, etc.).

In Figure 9, we can see different hyperspectral cubes extracted from the same bitstream with different spatial and spectral resolutions. The face of the cube is a color composition from different subbands. The spectral bands chosen for the color composition in the subresolution cube correspond to those of the original cube. Some slight differences from the original cube can be observed on the subresolution one, due to weighted averages from wavelet transform filtering spanning contiguous bands.

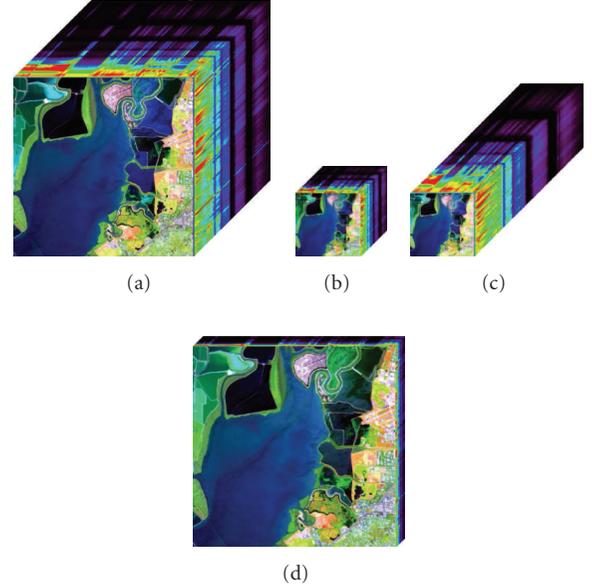


FIGURE 9: Example of hyperspectral cube with different spectral and spatial resolution decoded from the same bitstream. (a) is the original hyperspectral cube. (b) is 1/4 for spectral resolution and 1/4 for spatial resolution. (c) is full spectral resolution and 1/4 spatial resolution. (d) is full spatial resolution and 1/8 spectral resolution.

#### 6.4. ROI coding and selected decoding

The main interest of the present algorithm is in its flexibility. The bitstream obtained in the resolution scalable mode can be decoded at variable spectral and spatial resolutions for each data block. This is done reading, or transmitting, a minimum number of bits. Any area of the image can be decoded up to any spatial resolution, any spectral resolution and any bit plane. This property is illustrated in Figure 10. Most of the image background (area 1) is decoded at low spatial and spectral resolutions, dramatically reducing the amount of bits. Some specific areas are more detailed and, offer the full spectral resolution (area 2), the full spatial resolution (area 3), or both (area 4). The image from Figure 10 was obtained reading only 16907 bits from the 311598 bits belonging to the full codestream.

The region of interest can also be selected during the encoding while adjusting the number of bit planes to be encoded for a specific block. In the context of onboard processing, it would enable further reduction of the bit rate. The present encoder provides all these capabilities. For example, an external clouds detection loop could be added to adjust the compression parameter to reduce the resolution when clouds are detected. This would decrease the bit rate on these parts.

## 7. CONCLUSION

We have presented the 3D-SPIHT-RARS algorithm, an original extension of the 3D-SPIHT algorithm. This new algorithm enables resolution scalability for spatial and

TABLE 5: Bits read from full codestream for different resolution and quality for moffett3 image.

Resolution	Number of non decoded bit planes: 0				Number of non decoded bit planes: 2				Number of non decoded bit planes: 4			
	Full	1/2	1/4	1/8	Full	1/2	1/4	1/8	Full	1/2	1/4	1/8
bpppb read	5.309	1.569	0.247	0.038	2.857	0.989	0.198	0.033	1.016	0.475	0.132	0.027
RMSE	0.31	0.34	0.25	0.12	1.67	0.70	0.39	0.22	5.18	2.03	0.82	0.44
Time (s)	59.43	21.82	7.17	3.54	42.33	18.03	6.86	3.62	18.18	10.05	5.34	3.45
Resolution	Number of non decoded bit planes: 6				Number of non decoded bit planes: 8				Number of non decoded bit planes: 10			
	Full	1/2	1/4	1/8	Full	1/2	1/4	1/8	Full	1/2	1/4	1/8
bpppb read	0.327	0.203	0.079	0.020	0.104	0.077	0.039	0.013	0.030	0.025	0.016	0.007
RMSE	13.05	6.47	2.73	1.08	30.23	18.97	9.53	3.98	69.41	49.76	29.92	14.60
Time (s)	7.70	6.14	4.40	3.36	4.51	4.26	3.68	3.26	3.47	3.45	3.29	3.16

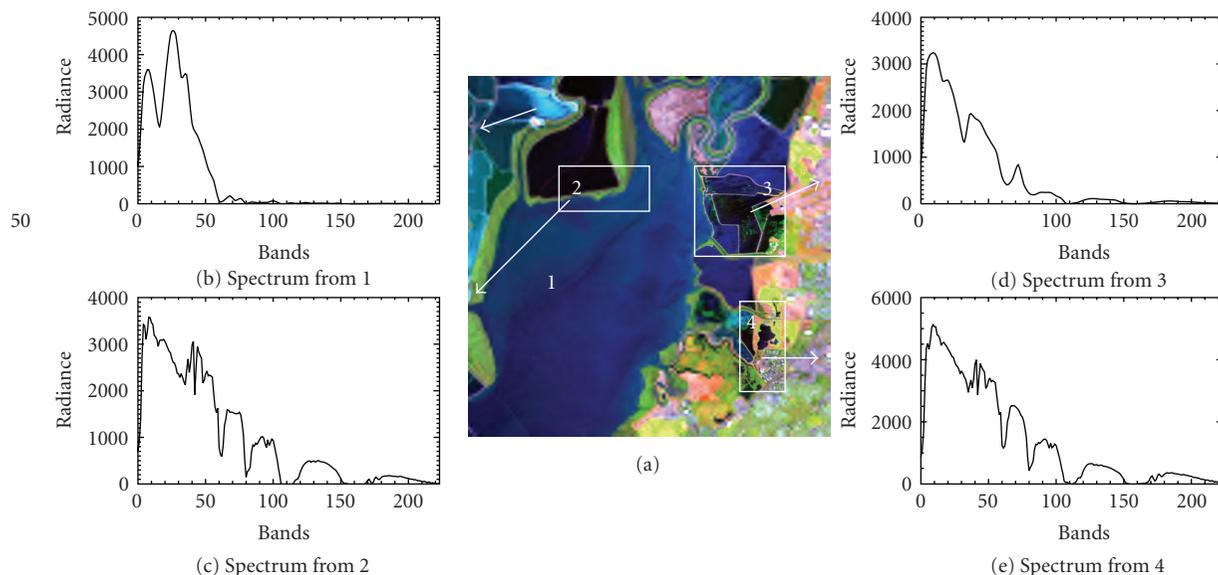


FIGURE 10: Example of a decompressed image with different spatial and spectral resolution for different areas. Background (area 1) is with low spatial resolution and low spectral resolution as is can be seen on the spectrum (b). Area 2 has low spatial resolution and high spectral resolution (c), area 3 has high spatial resolution, but low spectral resolution (d). Finally, area 4 has both high spectral and spatial resolutions. This decompressed image was obtained from a generic bitstream, reading the minimum amount of bits.

spectral dimensions independently and random access decoding. These properties are important to ease of access and processing of the data and were not introduced into SPIHT previously. Coding different areas of the image transform separately enable random access and region of interest coding with a reduction in memory usage during the compression. Furthermore, quality scalability for any resolution and area can be enabled by reorganization of the codestream. Thanks to the rate-distortion optimization between the different blocks, all these features are obtained without sacrificing compression performance. Most of these features seem also possible with the JPEG2000 standard. However, implementation providing multiresolution transforms is very recent and does not provide yet all the flexibility proposed here, particularly on the spectral direction. The granularity of the access is also finer with the proposed implementation.

The use of an arithmetic coder slightly increases compression performance, but at the cost of an increase in the complexity. It has to be highlighted that the 3D-SPIHT-RARS algorithm does not need to rely on arithmetic coding to obtain competitive results to JPEG2000.

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## Special Issue on The Human Microbiome and Infectious Diseases: Beyond Koch

### Call for Papers

A century after Robert Koch linked individual cultured microbes to specific diseases (Koch's postulates), it is increasingly apparent that the complex community of microorganisms associated with the human body (the "microbiome") plays a key role in health and disease. The National Institute of Health (NIH) recently announced the Human Microbiome Project and among its goals is to understand the relationship between host-associated microbial communities and disease. Many physicians and researchers, however, have only passing familiarity with the concepts involved in the study and therapeutic manipulation of complex microbial communities. The aims of this special issue are (1) to familiarize the readers with the concepts and methods for the study of complex microbial communities, (2) to demonstrate how changes in the indigenous microbial community can play a role in diseases such as antibiotic-associated diarrhea, bacterial vaginosis, and cystic fibrosis, and (3) to review how probiotics may hold promise for the therapeutic manipulation of the indigenous microbiota. Review articles and original research papers are being sought for this special issue.

Authors should follow the Interdisciplinary Perspectives on Infectious Diseases manuscript format described at the journal site <http://www.hindawi.com/journals/ipid/>. Prospective authors should submit an electronic copy of their complete manuscript through the journal Manuscript Tracking System at <http://mts.hindawi.com/>, according to the following timetable:

Manuscript Due	June 1, 2008
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### Guest Editors

**Vincent B. Young**, Division of Infectious Disease,  
Department of Internal Medicine, University of Michigan,  
Ann Arbor, MI 48109-5378, USA; [youngvi@umich.edu](mailto:youngvi@umich.edu)

**Robert A. Britton**, Department of Microbiology and  
Molecular Genetics, Michigan State University, East  
Lansing, MI 48824-4320, USA; [rbritton@msu.edu](mailto:rbritton@msu.edu)

**Thomas M. Schmidt**, Department of Microbiology and  
Molecular Genetics, Michigan State University, East  
Lansing, MI 48824-4320, USA; [tschmidt@msu.edu](mailto:tschmidt@msu.edu)

## Special Issue on Fairness in Radio Resource Management for Wireless Networks

### Call for Papers

Radio resource management (RRM) techniques (such as admission control, scheduling, sub-carrier allocation, channel assignment, power allocation, and rate control) are essential for maximizing the resource utilization and providing quality of service (QoS) in wireless networks.

In many cases, the performance metrics (e.g., overall throughput) can be optimized if opportunistic algorithms are employed. However, opportunistic RRM techniques always favor advantaged users who have good channel conditions and/or low interference levels. The problem becomes even worse when the wireless terminals have low mobility since the channel conditions become slowly varying (or even static), which might lead to long-term unfairness. The problem of fair resource allocation is more challenging in multi-hop wireless networks (e.g., mesh and multihop cellular networks).

The performance fairness can be introduced as one of the QoS requirements (e.g., as a condition on the minimum throughput per user). Fair RRM schemes might penalize advantaged users; hence, there should be a tradeoff between the overall system performance and the fairness requirements.

We are soliciting high-quality unpublished research papers addressing the problem of fairness of RRM techniques in wireless communication systems. Topics include (but are not limited to):

- Fairness of scheduling schemes in wireless networks
- Tradeoff between maximizing the overall throughput and achieving throughput fairness
- RRM fairness: problem definition and solution techniques
- Fairness performance in emerging wireless systems (WiMAX, ad hoc networks, mesh networks, etc.)
- Cross-layer RRM design with fairness
- Short-term and long-term fairness requirements
- Adaptive RRM to support fairness
- Fairness in cooperative wireless communications
- Issues and approaches for achieving fairness in multi-hop wireless networks

- RRM framework and QoS architecture
- Complexity and scalability issues
- Experimental and implementation results and issues
- Fairness in multiple-antenna transmission/reception systems
- Fairness in end-to-end QoS provisioning

Authors should follow the EURASIP Journal on Wireless Communications and Networking manuscript format described at the journal site <http://www.hindawi.com/journals/wcn/>. Prospective authors should submit an electronic copy of their complete manuscript through the journal Manuscript Tracking System at <http://mts.hindawi.com/>, according to the following timetable:

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### Guest Editors

**Mohamed Hossam Ahmed**, Faculty of Engineering and Applied Science, Memorial University of Newfoundland, St. John's, NF, Canada A1B 3V6; [mhahmed@engr.mun.ca](mailto:mhahmed@engr.mun.ca)

**Alagan Anpalagan**, Department of Electrical and Computer Engineering, Ryerson University, Toronto, ON, Canada M5G 2C5; [alagan@ee.ryerson.ca](mailto:alagan@ee.ryerson.ca)

**Kwang-Cheng Chen**, Department of Electrical Engineering, National Taiwan University, Taipei, Taiwan; [chenkc@cc.ee.ntu.edu.tw](mailto:chenkc@cc.ee.ntu.edu.tw)

**Zhu Han**, Department of Electrical and Computer Engineering, College of Engineering, Boise State University Boise, Idaho, USA; [zhuhan@boisestate.edu](mailto:zhuhan@boisestate.edu)

**Ekram Hossain**, Department of Electrical and Computer Engineering, University of Manitoba, Winnipeg, MB, Canada R3T 2N2; [ekram@ee.umanitoba.ca](mailto:ekram@ee.umanitoba.ca)

## Special Issue on Distributed Video Coding

### Call for Papers

Distributed source coding (DSC) is a new paradigm based on two information theory theorems: Slepian-Wolf and Wyner-Ziv. Basically, the Slepian-Wolf theorem states that, in the lossless case, the optimal rate achieved when performing joint encoding and decoding of two or more correlated sources can theoretically be reached by doing separate encoding and joint decoding. The Wyner-Ziv theorem extends this result to lossy coding. Based on this paradigm, a new video coding model is defined, referred to as distributed video coding (DVC), which relies on a new statistical framework, instead of the deterministic approach of conventional coding techniques such as MPEG standards.

DVC offers a number of potential advantages. It first allows for a flexible partitioning of the complexity between the encoder and decoder. Furthermore, due to its intrinsic joint source-channel coding framework, DVC is robust to channel errors. Because it does no longer rely on a prediction loop, DVC provides codec independent scalability. Finally, DVC is well suited for multiview coding by exploiting correlation between views without requiring communications between the cameras.

High-quality original papers are solicited for this special issue. Topics of interest include (but are not limited to):

- Architecture of DVC codec
- Coding efficiency improvement
- Side information generation
- Channel statistical modeling and channel coding
- Joint source-channel coding
- DVC for error resilience
- DVC-based scalable coding
- Multiview DVC
- Complexity analysis and reduction
- DSC principles applied to other applications such as encryption, authentication, biometrics, device forensics, query, and retrieval

Authors should follow the EURASIP Journal on Image and Video Processing manuscript format described at <http://www.hindawi.com/journals/ivp/>. Prospective authors should submit an electronic copy of their complete

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### Guest Editors

**Frederic Dufaux**, Ecole Polytechnique Fédérale de Lausanne, Lausanne, Switzerland; [frederic.dufaux@epfl.ch](mailto:frederic.dufaux@epfl.ch)

**Wen Gao**, School of Electronic Engineering and Computer Science, Peking University, Beijing, China; [wgao@pku.edu.cn](mailto:wgao@pku.edu.cn)

**Stefano Tubaro**, Dipartimento di Elettronica e Informazione, Politecnico di Milano, Milano, Italy; [stefano.tubaro@polimi.it](mailto:stefano.tubaro@polimi.it)

**Anthony Vetro**, Mitsubishi Electric Research Laboratories, Cambridge, MA, USA; [avetro@merl.com](mailto:avetro@merl.com)

## Special Issue on Broadcasting and Telecommunications at Crossroads: Impacts of NGN and Web 2.0 Concepts on the Future of IPTV

### Call for Papers

Broadcasting, telecommunications, and the Internet are converging. As transporting the traditional broadcast TV just over the Internet is probably not generating significant benefits for the customers, IPTV contemporary deployments provide additional capabilities, such as video-on-demand, network-based personal video recorders, time-shifted TV, and so forth. However, key values of IPTV probably originate from the availability of back control channel, allowing for interactive TV, as well as the integration with emerging multimedia group communication services, such as buddy lists featuring presence information and instant messaging, voice over IP, emailing, video conferencing and many more, enabling a new community TV experience. In this context, the notion of *IP multimedia subsystem- (IMS-) based IPTV* is currently gaining momentum in the international NGN and IPTV standardization, by providing the middleware to glue network intelligence and directly to the home or on the move, that is, wire or wireless, respectively, TV service provision.

In light of *Web 2.0*, The internet has become a user-generated multimedia content distribution medium for many different communities, including different types of video services: for example, YouTube, Apple Podcasts, or P2P video architectures like Joost. As the Internet is currently challenging many business cases for the provision of NGN communication services, it will also do so in the field of broadcasting and TV. There is no doubt that in the future NGN- and Web 2.0-enabled IPTV environments will be in a position to provide the required flexibility for supporting adequately the converged markets.

Authors should follow the International Journal of Digital Multimedia Broadcasting manuscript format described at the journal site <http://www.hindawi.com/journals/ijdmb/>. Prospective authors should submit an electronic copy of their complete manuscript through the journal Manuscript Tracking System at <http://mts.hindawi.com/> according to the following timetable:

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### Guest Editors

**Thomas Magedanz**, Fraunhofer Institute for Open Communication Systems (FOKUS), Kaiserin-Augusta-Allee 31, 10589 Berlin, Germany;  
[thomas.magedanz@fokus.fraunhofer.de](mailto:thomas.magedanz@fokus.fraunhofer.de)

**Heinrich Arnold**, Innovation Development Laboratory, Deutsche Telekom, Ernst-Reuter-Platz 7, 10587 Berlin, Germany; [heinrich.arnold@telekom.de](mailto:heinrich.arnold@telekom.de)

**Pete Distler**, Sprint Mobile Media Network, 6220 Sprint Parkway, Overland Park, KS 66251, USA;  
[peter.distler@sprint.com](mailto:peter.distler@sprint.com)

**Ravi S. Sharma**, Wee Kim Wee School of Communication and Information, Nanyang Technological University, 31 Nanyang Link, Singapore 637718; [asrsharma@ntu.edu.sg](mailto:asrsharma@ntu.edu.sg)

## Special Issue on OFDMA Architectures, Protocols, and Applications

### Call for Papers

Orthogonal frequency-division multiple access (OFDMA) technologies are currently attracting intensive attention in wireless communications to meet the ever-increasing demands arising from the explosive growth of Internet, multimedia, and broadband services. OFDMA-based systems are able to deliver high data rate, operate in the hostile multipath radio environment, and allow efficient sharing of limited resources such as spectrum and transmit power between multiple users. OFDMA has been used in the mobility mode of IEEE 802.16 WiMAX, is currently a working specification in 3GPP Long Term Evolution downlink, and is the candidate access method for the IEEE 802.22 “wireless regional area networks.” Clearly, recent advances in wireless communication technology have led to significant innovations that enable OFDMA-based wireless access networks to provide better quality-of-service (QoS) than ever with convenient and inexpensive deployment and mobility.

However, regardless of the technology used, OFDMA networks must not only be able to provide reliable and high quality broadband services, but also be implemented cost-effectively and be operated efficiently. OFDMA presents many of the advantages and challenges of OFDM systems for single users, and the extension to multiple users introduces many further challenges and opportunities, both on the physical layer and at higher layers. These requirements present many challenges in the design of network architectures and protocols, which have motivated a significant amount of research in the area. Also, many critical problems associated with the applications of OFDMA technologies in future wireless systems are still looking for efficient solutions. The aim of this special issue is to present a collection of high-quality research papers that report the latest research advances in this field from physical and network layers to practical applications. Original papers are solicited in all aspects of OFDMA techniques including physical layer issues, architectures, protocol designs, enabling technologies, theoretical studies, practical applications, and experimental prototypes. Topics of interest include, but are not limited to:

- Adaptive coding and modulation
- Signal processing for OFDMA
- Interference control techniques
- Bandwidth and resources allocation

- Efficient MAC protocol development
- Routing algorithms and congestion control schemes
- MAC and network layer management
- Cross-layer design and optimization
- Cooperative and game theoretic analysis
- Quality of service provisioning
- Network modeling and performance analysis
- Security and privacy management
- Broadband Wireless Access
- Testbed, experiment, implementation, standards, and practical applications

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### Guest Editors

**Victor C.M. Leung**, The University of British Columbia, 2329 West Mall Vancouver, BC, Canada, V6T 1Z4; [vleung@ece.ubc.ca](mailto:vleung@ece.ubc.ca)

**Alister G. Burr**, Department of Electronics, University of York, York, YO10-5DD, UK; [alister@ohm.york.ac.uk](mailto:alister@ohm.york.ac.uk)

**Lingyang Song**, Wireless Group, Philips Research Laboratories, Cross Oak Lane, Redhill, Surrey RH1 5HA, UK; [lingyang.song@philips.com](mailto:lingyang.song@philips.com)

**Yan Zhang**, Simula Research Laboratory, 1325 Lysaker, Norway; [yanzhang@ieee.org](mailto:yanzhang@ieee.org)

**Thomas Michael Bohnert**, Siemens AG, 80312 Munich, Germany; [thomas.bohnert@siemens.com](mailto:thomas.bohnert@siemens.com)

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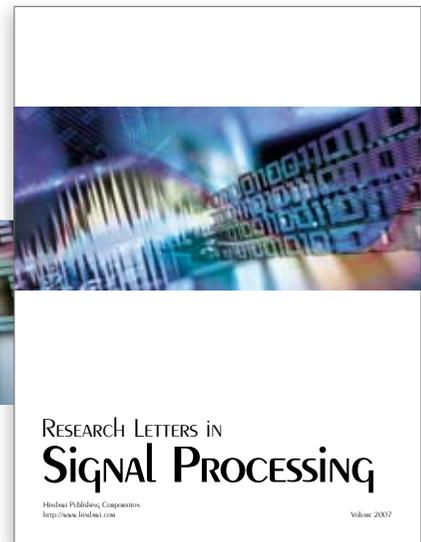
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