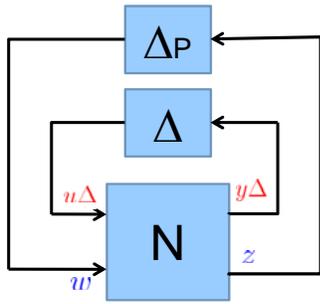


# Robust Controller Synthesis

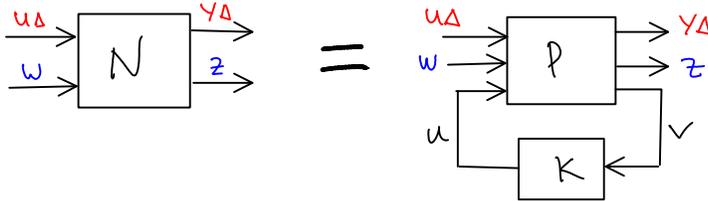
Requirement for Robust Performance (RP):  $N$  is internally stable (NS) and



has to be robustly stable for any allowed  $\Delta$ ,  $\|\Delta\|_\infty \leq 1$  and unstructured  $\Delta_p$ ,  $\|\Delta_p\|_\infty \leq 1$ .

From robust stability analysis, we know that this is equivalent to

$$\mu(N(j\omega)) < 1, \forall \omega$$



Therefore, in designing a robust controller, we want to achieve

- \*  $N$  is stable
- \*  $\mu(N(j\omega))$  is minimal (or at least less than 1)

Define:  $\hat{\Delta} = \begin{bmatrix} \Delta & 0 \\ 0 & \Delta_p \end{bmatrix}$ ,  $\mathcal{D}$  is the set of all matrices  $D$  such that  $D\hat{\Delta} = \hat{\Delta}D$

For each  $\omega$ ,  $\mu(N)$  can be upper bounded by

$$\mu(N) \leq \min_{D \in \mathcal{D}} \sigma_{\max}(DND^{-1}), \text{ which is a convex optimization problem}$$

Generally, for complex uncertainty with three or fewer blocks, this upper bound is tight.

## DK Iteration (8.12)

Generally, we are interested in designing a controller  $K$  according to the following minimization:

$$\min_{K, D \in \mathcal{D}} \sigma_{\max}(D N(K) D^{-1})$$

In practice, rather than solving this optimization for both  $D$  and  $K$  simultaneously we perform an iteration:

Step 0: Fix a  $D(\omega) \in \mathcal{D}$ , say  $D(\omega) = I$

Step 1: Obtain  $K$  by solving

$$\min_K \|D N(K) D^{-1}\|_{\infty} \quad \dots \quad (*)$$

Step 2: Using  $K$  from step 1, for every  $\omega$  find  $D(\omega)$  from

$$\min_{D(\omega) \in \mathcal{D}} \sigma_{\max}(D(\omega) N(j\omega) D^{-1}(\omega)) \quad \dots \quad (**)$$

Step 3: Fit the magnitude of  $D(\omega)$  to a stable and minimum phase transfer function  $D(s)$ , repeat step 1.

### Remarks:

- Step 1 is essentially an  $H_{\infty}$  synthesis with  $D$  as extra weight
- For each  $\omega$ , step 2 is a convex optimization problem
- The order of the controller in step 1 depends on the orders of the plant, weights, and  $D$ . For better performance,  $D(s)$  is usually high order, which leads to high order controller.

## Model Reduction (Ch 11)

General idea: Suppose that we are given a very complex plant model,  $G(s)$ . How can we come up with a simpler model  $G_a(s)$  that is still a good representation of the plant?

Complexity: The complexity of a linear <sup>proper</sup> model  $G(s)$  is typically measured as the minimum number of states required to represent it. This number is called the McMillan degree, which is also the number of system poles.

Reasoning: Computational: the number of states corresponds to the complexity of numerical simulation of the system. Also, many controller synthesis techniques use the state  $\rightarrow$  space representation

Physical: The McMillan degree is roughly related to the number of energy bearing variables in the system

Example:  $G(s) = \frac{1}{s} \begin{bmatrix} 1 & 0 \\ 1 & 2 \end{bmatrix}$

has pole polynomial  $\phi(s) = s^2$ , thus this is a second order system

A minimal realization:

$$\begin{aligned} \dot{x}_1 &= u_1 & \dot{x}_2 &= 2u_2 & \rightarrow A=0 & ; B &= \begin{bmatrix} 1 & 0 \\ 0 & 2 \end{bmatrix} & ; C &= \begin{bmatrix} 1 & 0 \\ 1 & 2 \end{bmatrix} \\ y_1 &= x_1 & y_2 &= x_1 + x_2 \end{aligned}$$

Compare with  $G(s) = \frac{1}{s} \begin{bmatrix} 1 & 0 \\ 1 & s \end{bmatrix}$ , whose pole polynomial is  $\phi(s) = s$

A minimal representation:

$$\left. \begin{aligned} \dot{x}_1 &= u_1 \\ y_1 &= x_1 \\ y_2 &= x_1 + u_2 \end{aligned} \right\} A=0 ; B=1 ; C = \begin{bmatrix} 1 \\ 1 \end{bmatrix} ; D = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

Quality of approximation is measured by the distance between  $G$  and  $G_a$ .  
In this course, we use  $H_{\infty}$  norm.

Thus, the general idea can be concretized as: Given a plant model with high order, how can we come up with a lower order / reduced order model that is close to the original model (in  $H_{\infty}$  sense)?